## ELLIOTT

## $\square$

| Volume | 2: | PROGRAMMING INFORMATION |
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| Part | 2: | PROGRAM DESCRIPTIONS |
| Section | 8: | QF (FLOATING POINT SUBROUTINES) |

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## Chapter 1: DESCRIPTION

### 1.1 INTRODUCTION

### 1.1.1 Purpose.

QF is used to perform operations on floatingpoint numbers.

### 1.1.2 Summary. <br> QF contains routines for operations

 corresponding to all the fixed-point operations except the function 15 (Functions 3 and 6 have special meanings in $Q F$ ).When entered, QF proceeds to interpret the instructions in the locations immediately following the entry-point in the user's program. Thus, operations are performed on floating-point numbers by placing in store the corresponding fixed-point instructions.

Two formats are available for floating-point numbers (see Paragraph 1.2.2).

### 1.1.3 Form of Distribution.

The program is distributed as an Elliott SIR mnemonic tape which has $Q F$ declared as a global identifier.
1.1.4 Method of Use.

See Paragraph 1.1.2. QF may be run in any program level.
1.1.5 Restrictions.

See Paragraph 1.2.4.
1.1.6 Accuracy.

If the result of the operation is $y$, multiplication gives a maximum error of $2^{-34} y$, division gives a maximum error of $2^{-32} y$. All other operations give a maximum error of $2^{-35} y$.

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## 1. 2 FUNCTIONS

| 1.2.1 $\quad$ Notation. |  |
| :--- | :--- |
| $\mathrm{x}(\mathrm{man})=$ | mantissa of floating-point number, x. |
| $\mathrm{x}(\mathrm{exp})=$ | exponent of floating-point number, x. |
| x | $=$ |
|  | a floating point number held in 2 or 3 words |
|  | from location X. |

### 1.2.2 Format.

One of two formats may be used to hold a floating-point number in store. Normally the packed format is used, but the unpacked format allows a wider range of numbers and slightly greater accuracy. The two formats are summarised in the table below. In the unpacked format, the mantissa is a double-length fraction held in two consecutive locations and the exponent is a single-length integer held in the next location. For the packed format, the mantissa is truncated and the exponent held in the seven least significant bits of the second store location. In this case the exponent must be in the range -64 to +63 .

| Format | Location | Bit 18 | Bits 17-8 |  |
| :--- | :---: | :---: | :---: | :---: |
| Packed | X | sign | most significant bits of <br> mantissa |  |
|  | $\mathrm{X}+1$ | 0 | least sig. <br> bits of <br> mantissa | exponent |
|  | $\mathrm{X}+1$ | 0 | most significant bits of <br> mantissa |  |
|  |  |  |  |  |

Negative number representation for exponent and mantissa is by the usual 2 's complement notation.

All internal working of QF uses the unpacked format.
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Examples of floating point numbers in the two formats:

| Number | Locn | Packed | Unpacked |
| :---: | :---: | :---: | :---: |
| 0.25= | X | 010000000000000000 | 010000000000000000 |
| 0. $5 \times 2^{-1}$ | X+1 | 000000000001111111 | 000000000000000000 |
|  | $\mathrm{X}+2$ | Not Used | 111111111111111111 |
| $\begin{aligned} & 1-2^{-37} \times 2^{63} \\ & \div 9.2 \times 10^{18} \end{aligned}$ | X | 011111111111111111 | 011111111111111111 |
|  | $\mathrm{X}+1$ | 011111111110111111 | 011111111110000000 |
|  | $\mathrm{X}+2$ | Not Used | 000000000000111111 |
| $\begin{aligned} & -1.0 \times 2^{-64} \\ & \doteqdot-0.5 \times 10^{-21} \end{aligned}$ | X | 100000000000000000 | 100000000000000000 |
|  | $\mathrm{X}+1$ | 000000000001000000 | 000000000000000000 |
|  | $\mathrm{X}+2$ | Not Used | 111111111111000000 |

1.2.3 Entry and Exit.

Entry is made by
$\left.\begin{array}{rll} & \begin{array}{rl}11 & Q F \\ 8 & Q F \\ \text { or } & 11\end{array} \\ & 8 & Q F \\ & Q F+2\end{array}\right\}$ to use packed format

QF proceeds to intexpret and execute the pseudo-program using

|  | a pseudo-accumulator | (FPA) |
| :--- | :--- | :--- |
| a pseudo-B-register | (FBREG) |  |
| and | a pseudo-S-register | $(\Omega F)$ |

See Paragraph 1.2.4 for the effects of each function.

Exit is made by placing +0 in the location after the last instruction to be interpreted. Control is then transferred to the location following the zero location. The Machine Accumulator and B register are not preserved. The pseudo. Accumulator and $B$ regidrers are not affected by entry and exit.
1.2.4 Available Operations.

See Paragraph 1.2.1 for notation used.

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Table 1

| Pseudo <br> Instruction. | Name | New Contents of |  | Remarks |
| :---: | :---: | :---: | :---: | :---: |
|  |  | FPA | X |  |
| 0 X | Load pseudo <br> B-register | f | x | $\begin{aligned} & \mathrm{b}:=\mathrm{C}(\mathrm{X}) \\ & \text { Notes } 1 \text { See } \\ & \text { N } 3 \text { 3 } \end{aligned}$ |
| 1 X | Add | f+x | x |  |
| 2 X | Negate \& Add | $\mathrm{x}-\mathrm{f}$ | x |  |
| 3 X | Exchange | x | f | Not a basic 903 operation |
| 4 X | Load FPA | x | x |  |
| 5 X | Store FPA | f | f | See Note 2 |
| 6 N | Conversion Routines | f | x | See Table 2 |
| 7 X | $\begin{aligned} & \text { Jump if } f= \\ & \text { zero } \end{aligned}$ | f | x | See Note 3 |
| 8 X | Jump | f | x | See Note 3 |
| 9 X | Jump if $\mathrm{f}<0$ | f | x | See Note 3 |
| 10 X | Count in store | f | "x+1" | See Note 3 |
| 11 X | Store pseudo SCR | f | $s$ | See Note 3 |
| 12 X | Multiply | $\mathrm{f}^{+} \mathrm{x}$ | x |  |
| 13 X | Divide | $f / \mathrm{x}$ | x | See Note 4 |
| 14 N | * $2^{N}$ | f* $2^{\text {N }}$ | x | N<4096 See Note 5 |
| 14 N | $\cdots 2^{N-8.192}$ | f* $2^{\text {N-8192 }}$ | x | $N \geq 4096$ |
| 15 N | error | f | x | See Paragraph 1.3 |

Notes
(1) The instruction 00 is interpreted as a terminator for floatingpoint working (See Paragraph 1.2.3).
(2) If packed format is in force during interpretation of a 3 or 5 instruction, then a test is made whether

$$
-64 \leq f(\exp )<+64
$$

If $f(\exp )<-64$ then $x:=0$ and the next instruction is interpreted.
If $f(\exp ) \geq+64$ then an error indication is output (See Paragraph 1.3).
(3) - These instructions operate on single word items. The instructions 7, 8 \& 9 may jump to another interpreted instruction: they must not jump out of the interpreted program except via a terminator.
(4) If an attempt is made to divide by zero an error indication is output (See Paragraph 1.3).
(5) The results of the following instructions are always standardised:

$$
1,2,12,13,14
$$

The instruction 140 may be used to standardise the contents of the FPA.
(6) Modified instructions may be used and, if they are, the contents of the pseudo B-register are added to the address digits before obeying an instruction.
(7) Literal addresses may not be used i. e. constants must be stored in the correct format by the user.

Table 2
The address of a 6 instruction determines its meaning.

| Function | Effect |
| :---: | :---: |
| 61 | Instructions interpreted after this assume packed format. |
| 62 | " " " " " unpacked " |
| 63 | Convert a single-1ength integer to a floating-point number and place the result in the FPA Location $Q F+3$ contains the address where the integer is held. |
| 64 | Form the integral part of the number in the FPA. Location $Q F+4$ contains the absolute address where the (single-length) integer is to be placed. This routine always rounds down. |
| 65 | Convert a fixed-point fraction to a floating-point number and place the result in the FPA, Location QF +5 contains the absolute address of the (single-length)fraction. |
| 66 | Convert the number in the FPA to a fixed-point fraction. Location $Q F+6$ contains the absolute address where the (single-1ength) fraction is to be placed. This routine always rounds down. |

The instructions 61 and 62 do not convert numbers; they define the action of following instructions:-

In 63 to 66 the arguments and the addresses in $Q F+3$ to QF+6 are unaffected. Error indications are output if an impermissible address is used or if overflow occurs. (See Paragraph 1. 3).

## Example

An integer is held in INT1 and a real number in RL2
The following section of program places the floating-point form of the first in RL1 and the entier of the second in INT2.
(SET ADDRESSES IN QF WORKSPACE)

| 4 | PSI1 | (PSI1 holds the address of INT1) |
| :--- | :--- | :--- |
| 5 | QF +3 |  |
| 4 | PSI2 | (PSI2 holds the address of INT2) |
| 5 | QF +4 |  |

(NOW PERFORM CONVERSIONS)

| 11 | QF | (ENTER QF) |
| :--- | :--- | :--- |
| 8 | QF +1 |  |
| 6 | 3 | (INT1 in the FPA) |
| 5 | RL1 |  |
| 4 | RL2 |  |
| 6 | 4 | (RL2 stored as integer) |
| +0 |  | (RETURN to normal working) |
| 8 | $;+0$ |  |

PSIl 0 INT1
PSI2 0 INT2

## 1. 3 ERROR INDICATIONS

## 1. 3. 1 <br> Standard Indication

If an error occurs a message is displayed output on the teleprinter followed by 5 inches of blank tape on the punch. Recovery can be made and the program continued by entry at 9 .
1.3.2 Errors Detected.

| Message | Significance | Effect of re-entry at 9 |
| :---: | :---: | :---: |
| QF! | Impermissible instruction | Dynamic stop at 9 |
| ROF! | Floating-point over-flow (5 or 13 instruction) | Result is taken as largest number of correct sign that can be held in packed format. |
| RTII: | Integer overflow | Result is taken as largest integer of correct sign. |
| RTF! | Fraction overflow ( 66 instruction) | Result is taken as largest fraction of correct sign. |

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### 1.4 METHOD USED

The following steps are carried out for each pseudoinstruction interpreted.
(1) The pseudo S-register is incremented.
(2) The function and the address (modified if required) bits of the interpreted instruction are stored.
(3) Control is transferred to the appropriate routine to execute the instruction.
(4) Control is returned to the interpreter (via a standardising routine for instructions $1,2,12,13,14$ ).

Subroutines from 903/Algol have been used for all arithmetic operations。

1. 5 STORE USED

The floating-point package occupies approximately
750
-900 locations.
1.6 TIME TAKEN

The following times are approximate:-

| Function Number | Operation | Times in Microseconds |
| :---: | :---: | :---: |
| 0 | Set pseudo B-register | 440 |
| 1 | Add | ```2150 to 3850 (average 2500- see Note 1)``` |
| 2 | Negate and Add | 2480 to 4180 (average 2800 see Note 1) |
| 3 | Exchange | 1980 packed or 1810 unpacked |
| 4 | Load FPA | 770 |
| 5 | Store FPA | 770 |
| 6 | Specifies format for the following instructions | 440 |
| 7 | Jump if zero | 440 |
| 8 | Jump | 440 |
| 9 | Jump if negative | 440 |
| 10 | Count in store | 400 |
| 11 | Store pseudo SCR | 440 |
| 12 | Multiply | $4300 \quad 2140 \quad 4700$ |
| 13 | Divide | $\$ 000$ to 10780 (average 9400 see Note 1) |
| 14 | Shift | 810 to 2510 (average 1210see Note 1) |
|  | ENTRY | 50 |
| 00 | EXIT | 150 |

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Notes (1) The time depends on the number of places the mantissa is shifted to standardise the result of the operation. The average given assumes a shift of 4 places.
(2) For modified instructions add $125 \mu \mathrm{~s}$ to the time taken.

